

Alphabets, Strings, and Languages

Key Concepts

Definition 1. An **alphabet** is a finite non-empty set. The elements of an alphabet are often called **symbols**.

Definition 2. A **string**, over an alphabet Σ , is a finite sequence of elements of Σ .

Definition 3. The **length** of a string ω , over an alphabet Σ , is the same as its length when it is considered as a sequence. The length of a string ω is often written as $|\omega|$.

- The empty string (that is, the string with length zero) will be denoted by λ .

Definition 4. For any alphabet Σ , Σ^* is the set of *all* strings over Σ .

Definition 5. Suppose that

$$\mu = a_1a_2 \dots a_n \quad \text{and} \quad \nu = b_1b_2 \dots b_m$$

are strings over an alphabet Σ (that is, suppose that $\mu, \nu \in \Sigma^*$) with lengths n and m , respectively, so that

$$\alpha_1, \alpha_2, \dots, \alpha_n, \beta_1, \beta_2, \dots, \beta_m \in \Sigma.$$

Then the **concatenation** of μ and ν , denoted by $\mu \cdot \nu$, is the string

$$\alpha_1\alpha_2 \dots \alpha_n\beta_1\beta_2 \dots \beta_m$$

over Σ , with length $n + m$, obtained by listing the symbols in ν after the symbols in μ .

Definition 6. If $\mu, \omega \in \Sigma^*$, for an alphabet Σ , then μ is a **substring** of ω if there exist strings $\nu, \varphi \in \Sigma^*$ such that $\omega = \nu \cdot \mu \cdot \varphi$.

Definition 7. Once again, suppose that $\mu, \omega \in \Sigma^*$, for an alphabet Σ . Then μ is a **prefix** of ω if there exists a string $\nu \in \Sigma^*$ such that

$$\omega = \mu \cdot \nu,$$

and μ is a **suffix** of ω if there exists a string $\nu \in \Sigma^*$ such that

$$\omega = \nu \cdot \mu.$$

Definition 8. A **language** over an alphabet Σ is a subset of Σ^* .