

Lecture #7: Regular Expressions and Regular Operations

What Will Happen During the Lecture

Review

The lecture presentation will begin with a **brief** review of the material in the preparatory video and documents for this lecture — and students will have the chance to ask questions about this.

Recognition and Application of a Regular Expression

Let $\Sigma = \{a, b, c\}$ — so that that does not include any of the symbols

$$\lambda, \emptyset, \Sigma, (,), \cup, \circ, *$$

and let

$$\Sigma_{\text{regex}} = \Sigma \cup \{\lambda, \emptyset, \text{"}\Sigma\text{"}, (,), \cup, \circ, *\} = \{a, b, c, \lambda, \emptyset, \text{"}\Sigma\text{"}, (,), \cup, \circ, *\}.$$

Given a string $\omega \in \Sigma_{\text{regex}}^*$, you might wish to do each of the following things:

- Decide whether ω is a regular expression over Σ .
- If ω is a regular expression over Σ , then — for a given string $\mu \in \Sigma^*$ — decide whether μ belongs to the language of ω .
- If ω is a regular expression over Σ , then *describe* the language of ω .

These problems will be discussed using the string

$$\omega = (((\Sigma)^* \circ a) \circ (\Sigma)^*) \circ (a \circ (\Sigma)^*) \in \Sigma_{\text{regex}}^*$$

and (for the second part of the problem) the strings $\mu_1 = abaca \in \Sigma^*$ and $\mu_2 = bac \in \Sigma^*$.

Designing a Regular Expression

If someone describes a language $L \subseteq \Sigma^*$ — and L is a regular language — how can you discover a regular expression ω (over the alphabet Σ) whose language is L ?

During the lecture presentation this question will be considered — as we design a regular expression — over the alphabet $\Sigma = \{a, b, c\}$, once again — for the language consisting of all strings in Σ^* that include an even number of copies of the symbol “a” — that is, for the language

$$L = \{\mu \in \Sigma^* \mid \text{the number of copies of “a” in } \mu \text{ is divisible by } 2\}.$$