

Analyzing the Running Time of a Simple Recursive Algorithm

Solutions for a Suggested Exercise

This exercise concerned the “Maximal Element in Part of an Integer Array” problem, considered in Reading #2, as well as the following algorithm — which was proved to correctly solve this problem in the exercise for that reading.

```
maxInRange2 ( integer[] A, integer low, integer high ) {  
  1. if (low == high) {  
  2.   return A[low]  
  } else {  
  3.   return max(maxInRange2(A, low, high - 1), A[high])  
  }  
}
```

1. You were first asked to use the uniform cost criterion to write a **recurrence** for the number $T_{\max}(k)$ of steps used by this algorithm when $0 \leq \text{low} \leq \text{high} \leq \text{A.length} - 1$ and $\text{high} - \text{low} + 1 = k \geq 1$.

Solution: Suppose first that $k = 1$. Then $\text{high} - \text{low} + 1 = k = 1$, so that $\text{low} = \text{high}$ and the test at line 1 is passed when checked (as part of an execution of the algorithm including these inputs). The algorithm then ends after a second step — at line 2 — is executed, so that $T_{\max}(1) = 2$.

Suppose, instead, that $k \geq 2$. Then $\text{high} - \text{low} + 1 = k \geq 2$, so that $\text{high} \geq \text{low} + 1$, and the test at line 1 fails when checked. In this case, execution continues, and ends, with the execution of the step at line 3.

However, this includes a recursive application of this algorithm with inputs $\text{low}' = \text{low}$ and $\text{high}' = \text{high} - 1$. Now

$$\begin{aligned} \text{high}' - \text{low}' + 1 &= (\text{high} - 1) - \text{low} + 1 \\ &= (\text{high} - \text{low} + 1) - 1 && \text{(reordering and grouping terms)} \\ &= k - 1 \end{aligned}$$

so that the recursive application of the algorithm at line 3 uses $T_{\max}(k - 1)$ steps.

Thus a **recurrence** for $T_{\max}(k)$ is as follows: If k is a positive integer then

$$T_{\max}(k) = \begin{cases} 2 & \text{if } k = 1, \\ T_{\max}(k - 1) + 2 & \text{if } k \geq 2. \end{cases}$$

2. You were next asked to guess a **solution** for this recurrence — that is, guess an expression for $T_{\max}(k)$ that is not in the form of a recurrence.

Solution: Notice that

- $T_{\max}(1) = 2$,
- $T_{\max}(2) = T_{\max}(1) + 2 = 2 + 2 = 4$,
- $T_{\max}(3) = T_{\max}(2) + 2 = 4 + 2 = 6$, and
- $T_{\max}(4) = T_{\max}(3) + 2 = 6 + 2 = 8$.

Thus after checking the value of $T_{\max}(k)$ for small positive values of k it might be reasonable to **guess** that $T_{\max}(k) = 2k$ for every positive integer k .

3. Finally you were asked to prove that your guess is correct.

Solution:

Claim: Suppose that $T_{\max} : \mathbb{N} \rightarrow \mathbb{N}$ such that, for every positive integer k ,

$$T_{\max}(k) = \begin{cases} 2 & \text{if } k = 1, \\ T_{\max}(k - 1) + 2 & \text{if } k \geq 2. \end{cases}$$

Then $T_{\max}(k) = 2k$ for every positive integer k .

Proof: By induction on k . The standard form of mathematical induction will be used.

Basis: If $k = 1$ then

$$\begin{aligned} T_{\max}(k) &= T_{\max}(1) \\ &= 2 && \text{(by the given recurrence for } T_{\max}(k)) \\ &= 2 \cdot 1 = 2k \end{aligned}$$

as required to establish the claim in this case.

Inductive Step: Let h be an integer such that $h \geq 1$. It is necessary and sufficient to use the following

Inductive Hypothesis: $T_{\max}(h) = 2h$.

to prove the following

Inductive Claim: $T_{\max}(h + 1) = 2(h + 1)$.

Since $h \geq 1$, $h + 1 \geq 2$. Therefore

$$\begin{aligned} T_{\max}(h + 1) &= T_{\max}((h + 1) - 1) + 2 && \text{(by the given recurrence for } T_{\max}(k)) \\ &= T_{\max}(h) + 2 \\ &= 2h + 2 && \text{(by the Inductive Hypothesis)} \\ &= 2(h + 1) \end{aligned}$$

— establishing the Inductive Claim, as needed to complete the Inductive Step and complete the proof of the claim. \square