An Energy Efficient Hierarchical Clustering Algorithm for Wireless Sensor Networks

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About the paper

- Appeared in INFOCOM 2003
- 11 pages
- 30 references
- Cited by 673 papers

Topic classification





Paper Structure





Related Work

- LCA
- LCA2
- WCA
- DCA
- DMAC
- LEACH
- non-purely hierarchical models

Assumptions (1/2)

- a) The sensors are distributed as per a homogeneous spatial Poisson process of intensity λ in 2-D space.
- b) All sensors transmit at the same power level so have the same radio range*r*.
- c)Data exchanged between two communicating sensors not within each others' radio range is forwarded by other sensors.

Assumptions (2/2)

- d) A distance of d between any sensor and its cluster head is equivalent to d/r hops.
- e) Each sensor uses one unit of energy to transmit or receive one unit of data.
- f) A routing infrastructure is in place;
 - so, only the sensors on the routing path forward the data.
- g) The communication environment is contention- and error-free. (No retransmit)

Single-level Clustering Algorithm

- Each sensor becomes a cluster-head (CH) with probability *p* and advertises itself as a CH to the sensors within its radio range.
- 2. The advertisement is forwarded to all the sensors that are no more than *k* hops away from CH. Any sensor that receives such advertisements and is not itself a CH joins the cluster of the closest CH.
- 3. Any sensor that is neither a CH nor has joined any cluster itself becomes a CH.



Optimal parameters

Probability of becoming a Clusterhead

$$p = \begin{bmatrix} \frac{1}{3c} + \frac{\sqrt[3]{2}}{3c(2+27c^2+3\sqrt{3}c\sqrt{27c^2+4})^{\frac{1}{3}}} \\ + \frac{(2+27c^2+3\sqrt{3}c\sqrt{27c^2+4})^{\frac{1}{3}}}{3c} \cdot \frac{1}{\sqrt[3]{2}} \end{bmatrix}^2$$

where $c = 3.06a\sqrt{\lambda}$.

Maximum number of hops allowed from a sensor to its clusterhead

$$k_1 = \left\lceil \frac{1}{r} \sqrt{\frac{-0.917 \ln(\alpha/7)}{p_1 \lambda}} \right\rceil$$



Results (1/4)





Results (2/4)



Results (3/4)

Number of Sensors (<i>n</i>)	Density (<i>d</i>)	Probability (p _{opt})	Maximum Number of Hops (k)
500	5	0.1012	5
1000	10	0.0792	4
1500	15	0.0688	3
2000	20	0.0622	3
2500	25	0.0576	3
3000	30	0.0541	3



Results (4/4)



Hierarchical Clustering Alg. (1/2)

- Each sensor decides to become a level-1 CH with p_1 and advertises itself as a CH to the sensors within its radio range.
- This advertisement is forwarded to all the sensors within k_1 hops of the advertising. Each sensor that receives an advertisement joins the cluster of the closest level-1 CH; the remaining sensors become forced level-1 CHs.

Hierarchical Clustering Alg. (2/2)

- Level-1 CHs then elect themselves as level-2 CHs with a certain probability p_2 and broadcast their decision of becoming a level-2 CH. This decision is forwarded to all the sensors within k_2 hops.
- The level-1 CHs that receive the advertisements from level-2 CHs joins the cluster of the closest level-2 CH. All other level-1 CHs become forced level-2 CHs.
- CH at level 3,4,...,h are chosen in similar fashion, with probabilities p_3 , p_4 ,..., p_h respectively

Optimal parameters

Probability of becoming a Clusterhead

$$\begin{split} E[C] &= E[E[C \mid N = n]] \\ &= \lambda A \prod_{i=1}^{h} p_i \left[\frac{0.765a}{r} \right] \\ &+ \lambda A \sum_{i=1}^{h} (1 - p_i) \prod_{j=1}^{i-1} \left(p_j \right) \left[\frac{1}{2r \sqrt{\lambda \prod_{j=1}^{i} p_j}} \right]. \end{split}$$

Maximum number of hops allowed from a sensor to its clusterhead

$$k_{i} = \begin{bmatrix} \frac{1}{r} \sqrt{\frac{-0.917 \ln(\alpha/7)}{\lambda \prod_{j=1}^{i} p_{j}}} \end{bmatrix}$$



Results (1/4)





Results (2/4)





Results (3/4)





Results (4/4)





Q&A



Conclusion (1/2)

- A distributed algorithm proposed for organizing sensors into a hierarchy of clusters.
- Objective is minimizing the total energy spent in the system to communicate the information gathered by these sensors to the information-processing center.
- The optimal parameter values found for these algorithms that minimize the energy spent in the network.

Conclusion (2/2)

- Assumed that the communication environment is contention and error free.
- In a contention-free environment, the algorithm has a time complexity of $O(k_1 + k_2 + ... + k_h)$, a significant improvement over the many O(n) clustering algorithms in the literature.
- The algorithm is suitable for networks of large number of nodes.



Future Work

 Consider an underlying medium access protocol and investigate how that would affect the optimal probabilities of becoming a CH and the run-time of the algorithm.



Thank You

