

Recursion for Data Structures

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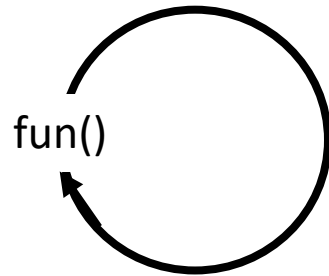
Recursion

- Definition:
 - See Recursion
 - Defining something in terms of itself
 - Generally using a smaller or simpler version
- Recursive Function
 - A function that calls itself

Recursion

- A programming technique whereby a function calls itself either directly or indirectly

```
def fun ():  
    :  
    fun ()  
    :
```

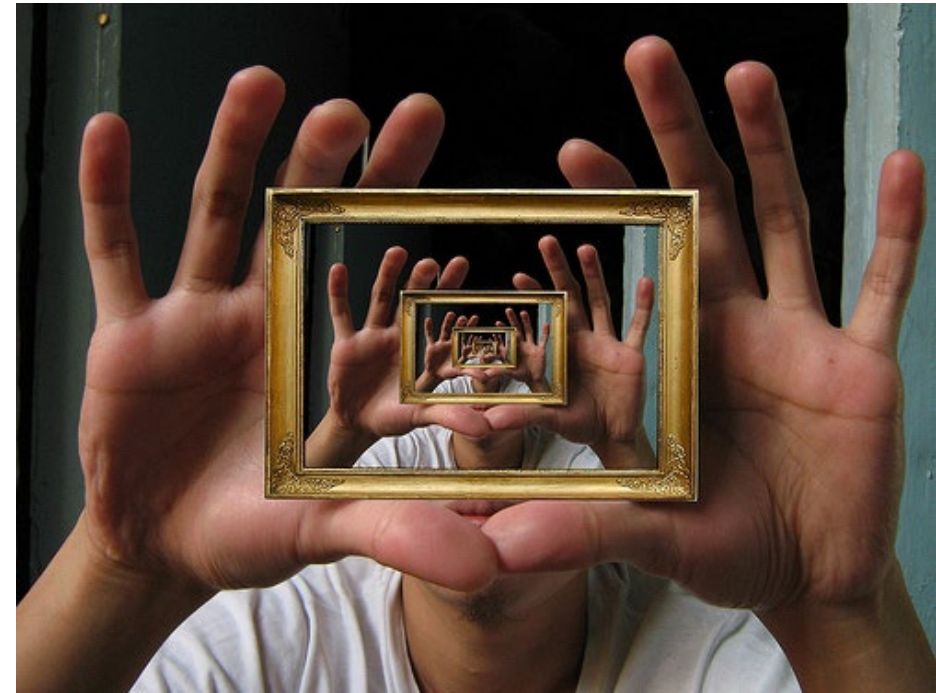


```
def fun1 ():  
    fun2 ()
```

```
def fun2 ():  
    fun1 ()
```

indirectly (mutual recursive)

Directly



Tail Recursion

- a tail call is a subroutine call performed as the final action of a function (Ex. A final return statement function call)
- If a tail call calls the function itself (as the last thing the function does), then this function is called tail recursive

#Not tail-recursive

```
def fact(n):
```

```
    if (n == 0):
```

```
        return 1
```

```
    return n * fact(n-1)
```

#Tail recursive

```
def fact(n, a = 1):
```

```
    if (n == 1):
```

```
        return a
```

```
    return fact(n - 1, n * a)
```

Tail Recursion

- Typically when a function is called a record of memory (stack frame) is stored to track state from the previous function call
- But a tail recursive function doesn't need anything from before the final function call
- Modern compilers can identify this and reduce the process into an iterative loop for you. Although this is not guaranteed.
 - Java is still a procedural language, there is a class of languages called functional languages which will often guarantee this behaviour
- **Neither Java or Python optimizes tail recursion**

Factorial

A Small Example - Factorial

Computing n factorial (n!):

- Defined as the result of multiplying all numbers from n to 1 for all $n > 0$. If $n == 0$, then result is 1.

$$n! = \begin{cases} 1 & n = 0 \\ n \times (n - 1)! & n > 0 \end{cases}$$

- $0! \rightarrow 1$
- $1! \rightarrow 1 * 1 = 1$
- $2! \rightarrow 2 * 1 * 1 = 2$
- $3! \rightarrow 3 * 2 * 1 * 1 = 6$

A Small Example

Recursive Definition!

- Another solution
 - BASE CASE,
 - RECURSIVE CASE,

$$0! == 1$$

$$n! == n * (n-1)!$$

A Small Example - Factorial

```
public static int factorial(int n){
    if(n < 0){
        throw new IllegalArgumentException("...");
    }
    if (n == 0){
        return 1;
    }
    return n * factorial(n-1);
}
```

A Small Example - Factorial

```
public static int factorial(int n){  
    if(n < 0){  
        throw new IllegalArgumentException("...");  
    }  
    if (n == 0){  
        return 1;  
    }  
    return n * factorial(n-1);  
}
```

We saw a BigInteger version in TicTacToe

Fibonacci

Fibonacci Numbers

- A sequence of values:
 - 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, ...
- Defined recursively:
 - By definition:
 - fib(0) is 0
 - fib(1) is 1
 - Remaining values:
 - $\text{fib}(n) = \text{fib}(n-1) + \text{fib}(n-2)$

Fibonacci Numbers

```
public int fibonacci(int n){
    if(n < 0){
        throw new IllegalArgumentException("...");
    }
    if(n == 0){
        return 0;
    }
    if(n == 1){
        return 1;
    }
    return fibonacci(n-1) + fibonacci(n-2);
}
```

Recursion – Beyond Algorithms

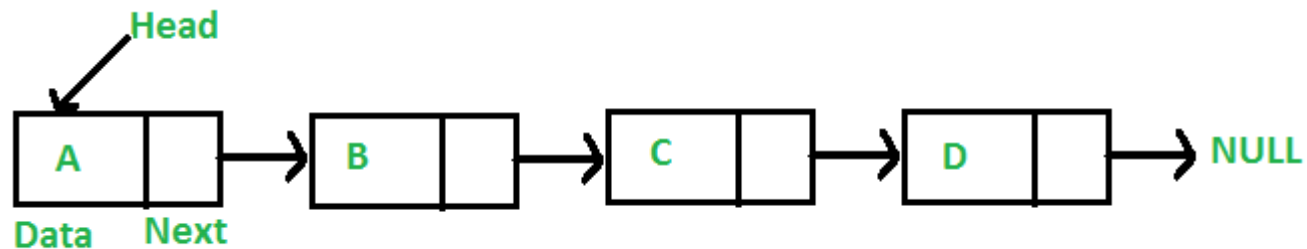
Not all about code

- We are going to make two recursive data structures
 1. A LinkedList
 2. A BinarySearchTree

Linked List

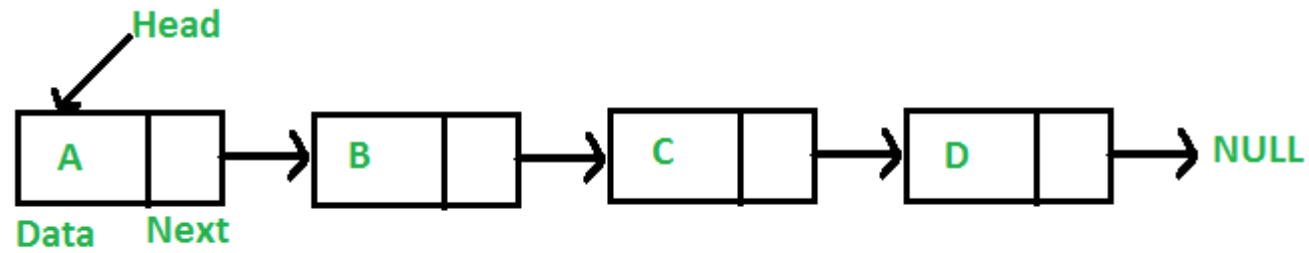
Linked List

- Data is stored structurally, 1 NODE per piece of data
- We have connection to start of list called HEAD
- Each NODE connects to another NODE to connect to more data, or NULL if there is no more data



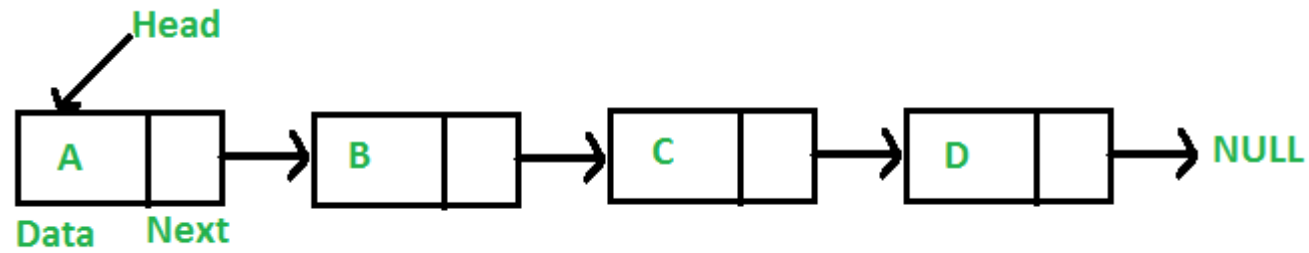
Linked List

- We want 2 classes
- One LinkedList and one for idea of Node



Linked List

- We want 2 classes
- One LinkedList and one for idea of Node

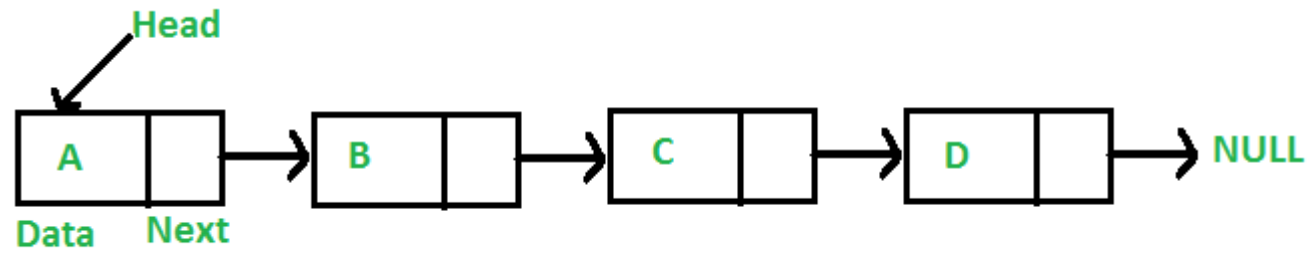


```
public class LinkedList {  
}
```

```
public class Node {  
}
```

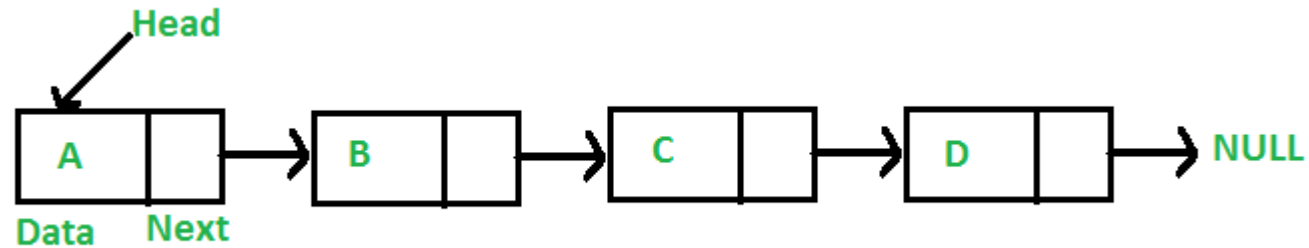
Linked List

- Does any other part of code need to know about a Node every?
- No! (Solution is internal class)



```
public class LinkedList {  
    private class Node {  
    }  
}
```

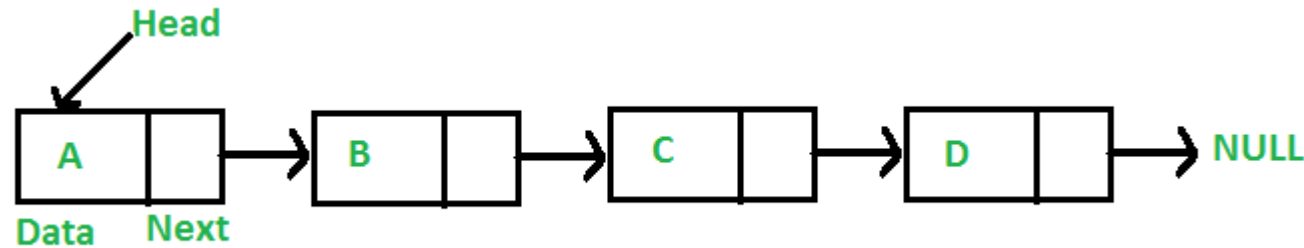
Linked List



```
public class LinkedList {  
    private Node head;  
    private class Node {  
        Object data;  
        Node next;  
    }  
}
```

- We don't have to worry the same amount about modifiers for internal private class as the data in an node is never exposed in the same way
- We can still make it private and make accessors but for simplicity we'll use Node more like just storage package of data

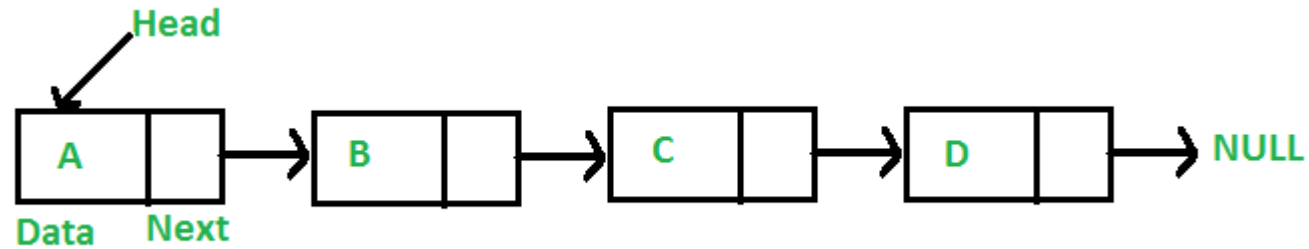
Linked List



```
public class LinkedList {  
    private Node head;  
    private class Node {  
        Object data;  
        Node next;  
    }  
}
```

- We don't have to worry the same amount about modifiers for internal private class as the data in an node is never exposed in the same way
- We can off course still make it private and make accessors but for simplicity we'll use Node more like just storage package of data

Linked List



```
public class LinkedList<E> {  
    private Node head;  
    private class Node {  
        E data;  
        Node next;  
    }  
}
```

- Generics!

Starting point

```
public class LinkedList<E> {  
    private Node head;  
    private class Node {  
        E data;  
        Node next;  
    }  
    public LinkedList(){  
        head = null;  
    }  
}
```


Make it a Java List interface type

```
import java.util.*;

public class LinkedList<E> implements List<E> {
    private Node head;
    private class Node {
        E data;
        Node next;
    }
    public LinkedList(){
        head = null;
    }
}
```

We'll stick with easiest to start (non index)

```
import java.util.*;
```

```
public class LinkedList<E> implements List<E> {  
    private Node head;  
    private class Node {  
        E data;  
        Node next;  
    }  
    public LinkedList(){  
        head = null;  
    }  
}
```

```
public int size()  
public boolean isEmpty()  
public boolean contains(Object o)  
public boolean add(E e)  
public boolean remove(Object o)  
public void clear()
```

We'll stick with easiest to start (non index)

```
import java.util.*;
```

```
public class LinkedList<E> implements List<E> {  
    private Node head;  
    private class Node {  
        E data;  
        Node next;  
    }  
    public LinkedList(){  
        head = null;  
    }  
}
```

```
public int size()  
public boolean isEmpty()  
public boolean contains(Object o)  
public boolean add(E e)  
public boolean remove(Object o)  
public void clear()
```

Size()

- Right now the only way to know this is to loop and count how many things are stored!

Size()

- Let's instead track a efficient integer for this value

```
public class LinkedList<E> implements List<E> {  
    private Node head;  
    private int size;  
    public LinkedList(){  
        head = null;  
        size = 0;  
    }  
    @Override  
    public int size() {  
        return size;  
    }  
}
```

```
@Override  
public void clear() {  
    head = null;  
    size = 0;  
}
```

```
@Override  
public boolean isEmpty() {  
    return size == 0;  
}
```

contains()

- Iterative contains and recursive contains

@Override

```
public boolean contains(Object o) {  
    Node curr = head;  
    while (curr != null) {  
        if (curr.data.equals(o)) {  
            return true;  
        }  
        curr = curr.next;  
    }  
    return false;  
}
```

@Override

```
public boolean contains(Object o) {  
    return recContains(head, o);  
}  
  
private boolean recContains(Node node, Object o) {  
    if (node == null) { //base case  
        return false;  
    }  
    if (node.data.equals(o)) { // base case  
        return true;  
    }  
    return recContains(node.next, o); // iterative case  
}
```

toString()

@Override

```
public String toString(){
    StringBuilder sb = new StringBuilder();
    sb.append("[");
    Node curr = head;
    while(curr != null){
        sb.append(curr.data);
        if(curr.next != null) {
            sb.append(", ");
        }
        curr = curr.next;
    }
    sb.append("]");
    return sb.toString();
}
```

- This is an iterative version of toString()
- We can at the same time program a recursive version

toString()

@Override

```
public String toString() {
    StringBuilder sb = new StringBuilder();
    sb.append("[");
    recToString(head, sb);
    sb.append("]");
    return sb.toString();
}

private void recToString(Node node, StringBuilder sb) {
    if (node == null) {
        return;
    }
    sb.append(node.data);
    if (node.next != null) {
        sb.append(", ");
    }
    recToString(node.next, sb);
}
```

- This is a recursive version of toString()
- Base case is we are at end and are done!
- Recursive case is we add current node data, and then go to next node
 - (there is a sub-case where we add the comma only if we aren't last entry (indicated by next being null))

add()

- Iterative add
- We make a new head node if currently empty list
- Otherwise we loop until end of list and then add a new node there
- Slow!!!!!!!
- Solutions?

@Override

```
public boolean add(E e) {  
    if(head == null){  
        head = new Node();  
        head.data = e;  
        size++;  
        return true;  
    }  
    Node curr = head;  
    while (curr != null) {  
        if (curr.next == null) {  
            curr.next = new Node();  
            curr.next.data = e;  
            size++;  
            return true;  
        }  
        curr = curr.next;  
    }  
    return false;  
}
```

add()

- Iterative add
- We make a new head node if currently empty list
- Otherwise we loop until end of list and then add a new node there
- Slow!!!!!!!
- Solutions? Track a pointer to end of list node!

@Override

```
public boolean add(E e) {  
    if(head == null){  
        head = new Node();  
        head.data = e;  
        size++;  
        return true;  
    }  
    Node curr = head;  
    while (curr != null) {  
        if (curr.next == null) {  
            curr.next = new Node();  
            curr.next.data = e;  
            size++;  
            return true;  
        }  
        curr = curr.next;  
    }  
    return false;  
}
```

add()

- Recursive add
- We make a new head node if currently empty list
- Otherwise, explore (recursive case) until end of list and then add (base case)

@Override

```
public boolean add(E e) {  
    if(head == null){  
        head = new Node();  
        head.data = e;  
        size++;  
        return true;  
    }  
    return recAdd(head, e);  
}
```

```
private boolean recAdd(Node node, E e) {  
    if(node.next == null){  
        node.next = new Node();  
        node.next.data = e;  
        size++;  
        return true;  
    }  
    return recAdd(node.next, e);  
}
```

remove()

- Iterative remove

```
@Override
public boolean remove(Object o) {
    Node curr = head;
    Node prev = null;
    while (curr != null) {
        //What if I found it?
        prev = curr;
        curr = curr.next;
    }
    return false;
}
```

remove()

- Iterative remove
- If found at head, then we set head to skip past
- If found later we set previous to skip past current to the next

```
@Override
public boolean remove(Object o) {
    Node curr = head;
    Node prev = null;
    while (curr != null) {
        if (curr.data.equals(o)) {
            if (curr == head) {
                head = head.next;
            } else {
                prev.next = curr.next;
            }
            size--;
            return true;
        }
        prev = curr;
        curr = curr.next;
    }
    return false;
}
```

remove()

- Recursive remove
- Deal with empty start of list
- Then deal with removing at start of list, otherwise enter recursion

- Recursion actually checks delete of next node, and deletes by skipping past it from current node
- If we reach end we have base case to return false

@Override

```
public boolean remove(Object o) {  
    if (head == null) {  
        return false;  
    }  
    if (head.data.equals(o)) {  
        head = head.next;  
        return true;  
    } else {  
        return recRemove(head, o);  
    }  
}
```

```
public boolean recRemove(Node node, Object o) {  
    if (node.next == null) {  
        return false;  
    }  
    if (node.next.data.equals(o)) {  
        node.next = node.next.next;  
        size--;  
        return true;  
    }  
    return false;  
}
```

LinkedList (Make it a List)

```
public static void main(String[] args) {  
    LinkedList<Integer> nal = new LinkedList <>();  
    nal.add(1);  
    nal.add(2);  
    nal.add(3);  
    System.out.println(nal);  
}
```

[1, 2, 3]

Named LinkedList (Make it a List)

```
public static void main(String[] args) {  
    LinkedList<Integer> nal = new LinkedList<>();  
    System.out.println(nal.size() + " " + nal.isEmpty());  
    nal.add(1);  
    System.out.println(nal.size() + " " + nal.isEmpty());  
    nal.add(2);  
    System.out.println(nal.size() + " " + nal.isEmpty());  
    nal.add(3);  
    System.out.println(nal.size() + " " + nal.isEmpty());  
}
```

0 true
1 false
2 false
3 false

Named LinkedList (Make it a List)

```
public static void main(String[] args) {  
    LinkedList<Integer> nal = new LinkedList<>();  
    nal.add(1);    nal.add(2);    nal.add(3);  
    System.out.println(nal);  
    System.out.println(nal.contains(1));  
    System.out.println(nal.contains(4));  
    nal.remove(new Integer(1));  
    System.out.println(nal);  
    System.out.println(nal.contains(1));  
    nal.clear();  
    System.out.println(nal);  
}
```

[1, 2, 3]

true

false

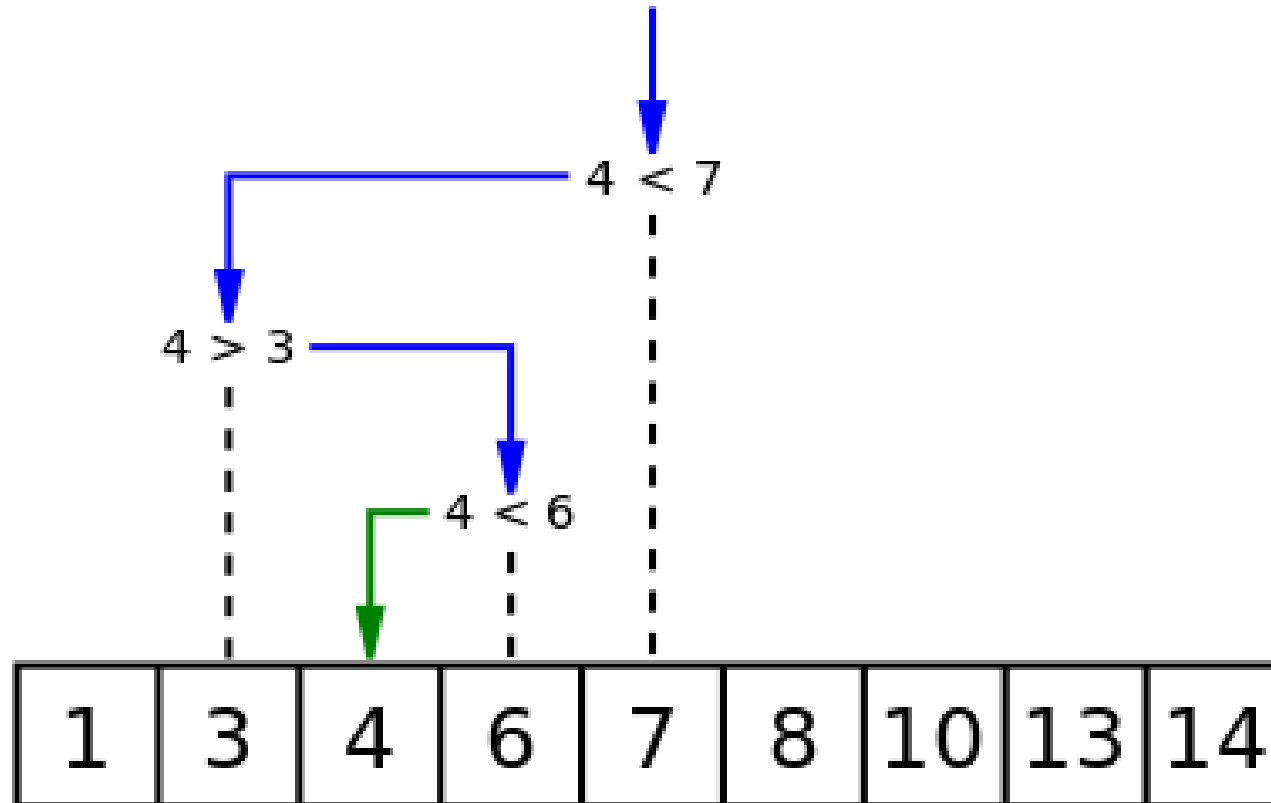
[2, 3]

false

[]

Binary Search

Binary Search



Binary Search

```
public static Integer binarySearch(Comparable[] values, Comparable search){
    return binarySearch(values, search, 0, values.length-1);
}
private static Integer binarySearch(Comparable[] values, Comparable search, int start, int end){
    if(start > end){
        return null;
    }
    int middle = (start + end) / 2;
    int compare = search.compareTo(values[middle]);
    if(compare == 0){
        return middle;
    }
    if (compare < 0){
        return binarySearch(values, search, start, middle-1);
    }else{
        return binarySearch(values, search, middle+1, end);
    }
}
```

Binary Search

```
Integer[] values = new Integer[]{0,2,4,6,8,10,12,14,16,18,20};
for (int i = -1; i <= 21; i++) {
    System.out.println(i+" -> "+binarySearch(values, new Integer(i)));
}
```

```
-1 -> null
0 -> 0
1 -> null
2 -> 1
3 -> null
4 -> 2
5 -> null
6 -> 3
7 -> null
8 -> 4
9 -> null
10 -> 5
11 -> null
12 -> 6
13 -> null
14 -> 7
15 -> null
16 -> 8
17 -> null
18 -> 9
19 -> null
20 -> 10
21 -> null
```

Binary Search

```
Integer[] values = new Integer[]{20,18,16,14,12,10,8,6,4,2,0};  
for (int i = -1; i <= 21; i++) {  
    System.out.println(i+" -> "+binarySearch(values, new Integer(i)));  
}
```

```
-1 -> null  
0 -> null  
1 -> null  
2 -> null  
3 -> null  
4 -> null  
5 -> null  
6 -> null  
7 -> null  
8 -> null  
9 -> null  
10 -> 5  
11 -> null  
12 -> null  
13 -> null  
14 -> null  
15 -> null  
16 -> null  
17 -> null  
18 -> null  
19 -> null  
20 -> null  
21 -> null
```

Binary Search

```
Integer[] values = new Integer[]{20,18,16,14,12,10,8,6,4,2,0};  
for (int i = -1; i <= 21; i++) {  
    System.out.println(i+" -> "+binarySearch(values, new Integer(i)));  
}
```

-1 -> null
0 -> null
1 -> null
2 -> null
3 -> null
4 -> null
5 -> null
6 -> null
7 -> null
8 -> null
9 -> null
10 -> 5
11 -> null
12 -> null
13 -> null
14 -> null
15 -> null
16 -> null
17 -> null
18 -> null
19 -> null
20 -> null
21 -> null

Binary Search

```
Integer[] values = new Integer[]{20,18,16,14,12,10,8,6,4,2,0};  
Arrays.sort(values);  
for (int i = -1; i <= 21; i++) {  
    System.out.println(i+" -> "+binarySearch(values, new Integer(i)));  
}
```

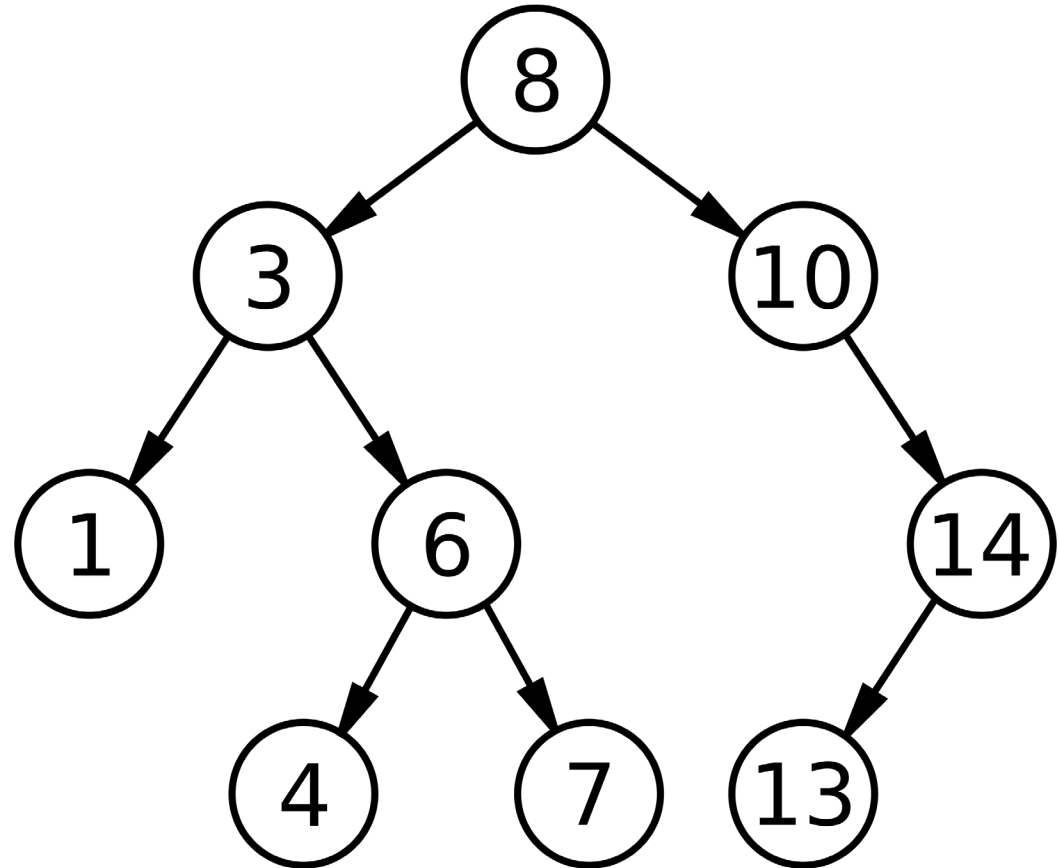
```
-1 -> null  
0 -> 0  
1 -> null  
2 -> 1  
3 -> null  
4 -> 2  
5 -> null  
6 -> 3  
7 -> null  
8 -> 4  
9 -> null  
10 -> 5  
11 -> null  
12 -> 6  
13 -> null  
14 -> 7  
15 -> null  
16 -> 8  
17 -> null  
18 -> 9  
19 -> null  
20 -> 10  
21 -> null
```


Binary Search Tree

Binary Search Tree

- Data on left is always less than data in middle
- Data on right is always more than data in middle

- This is true for every node



Binary Search Tree

- Data on left is always less than data in middle
- Data on right is always more than data in middle

- This is true for every node

```
public class BinarySearchTree<K extends Comparable> {  
    private Node root;  
    private class Node{  
        Comparable<K> key;  
        Node left, right;  
    }  
}
```

add

- If empty then adding is simple otherwise we have to make choices of middle, left, or right

```
public boolean add(K key){  
    if(root == null){  
        root = new Node();  
        root.key = key;  
        return true;  
    }  
    return recAdd(root, key);  
}
```

add

```
private boolean recAdd(Node node, K key) {
    int compare = key.compareTo(node.key);
    if(compare == 0){
        return false;
    }
    if(compare < 0) {
        if (node.left == null) {
            node.left = new Node();
            node.left.key = key;
            return true;
        }
        return recAdd(node.left, key);
    }
}
```

```
else{
    if (node.right == null) {
        node.right = new Node();
        node.right.key = key;
        return true;
    }
    return recAdd(node.right, key);
}
}
```

contains

- We start at root,
- Either we are at a null, find it, or go left, or right

```
public boolean contains(K key){  
    return recContains(root, key);  
}
```

```
private boolean recContains(Node node, K key) {  
    if(node == null){  
        return false;  
    }  
    int compare = key.compareTo(node.key);  
    if(compare == 0){  
        return true;  
    }  
    if(compare < 0) {  
        return recContains(node.left, key);  
    }else{  
        return recContains(node.right, key);  
    }  
}
```

toString

- We build outer part, then we go through tree, left before current, current, then right
- This gets us sorted order

@Override

```
public String toString(){
    StringBuilder sb = new StringBuilder();
    sb.append("[");
    recToString(root, sb);
    sb.delete(sb.length()-2,sb.length());
    sb.append("]");
    return sb.toString();
}
```

```
private void recToString(Node node, StringBuilder sb) {
    if(node == null){
        return;
    }
    recToString(node.left, sb);
    sb.append(node.key);
    sb.append(", ");
    recToString(node.right, sb);
}
```

toString

```
Integer[] values = new Integer[]{0,2,4,6,8,10,12,14,16,18,20};
BinarySearchTree<Integer> bst = new BinarySearchTree<>();
for(Integer value: values){
    bst.add(value);
}
System.out.println(bst);
for (int i = -1; i <= 21; i++) {
    System.out.println(i+" -> "+bst.contains(new Integer(i)));
}
```

```
[0, 2, 4, 6, 8, 10, 12, 14, 16, 18, 20]
-1 -> false
0 -> true
1 -> false
2 -> true
3 -> false
4 -> true
5 -> false
6 -> true
7 -> false
8 -> true
9 -> false
10 -> true
11 -> false
12 -> true
13 -> false
14 -> true
15 -> false
16 -> true
17 -> false
18 -> true
19 -> false
20 -> true
21 -> false
```

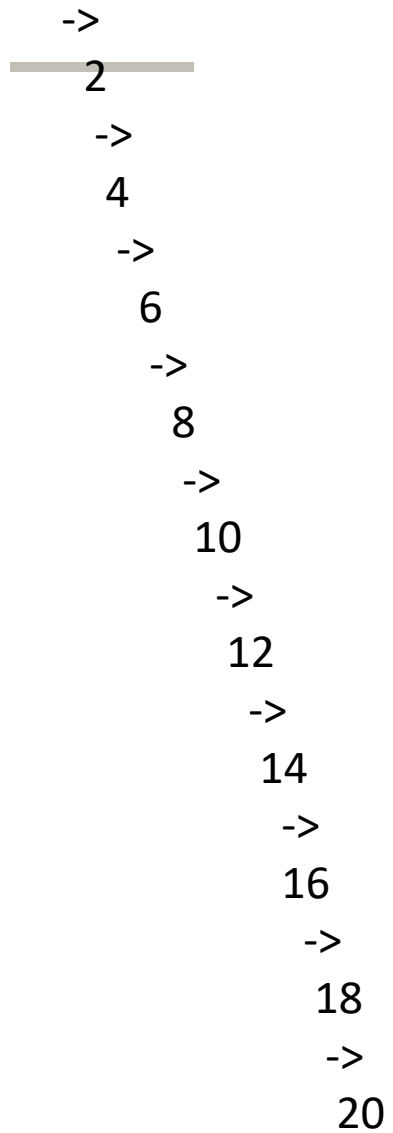

toString

```
Integer[] values = new Integer[]{20,18,16,14,12,10,8,6,4,2,0};
BinarySearchTree<Integer> bst = new BinarySearchTree<>();
for(Integer value: values){
    bst.add(value);
}
System.out.println(bst);
for (int i = -1; i <= 21; i++) {
    System.out.println(i+" -> "+bst.contains(new Integer(i)));
}
```

```
[0, 2, 4, 6, 8, 10, 12, 14, 16, 18, 20]
-1 -> false
0 -> true
1 -> false
2 -> true
3 -> false
4 -> true
5 -> false
6 -> true
7 -> false
8 -> true
9 -> false
10 -> true
11 -> false
12 -> true
13 -> false
14 -> true
15 -> false
16 -> true
17 -> false
18 -> true
19 -> false
20 -> true
21 -> false
```

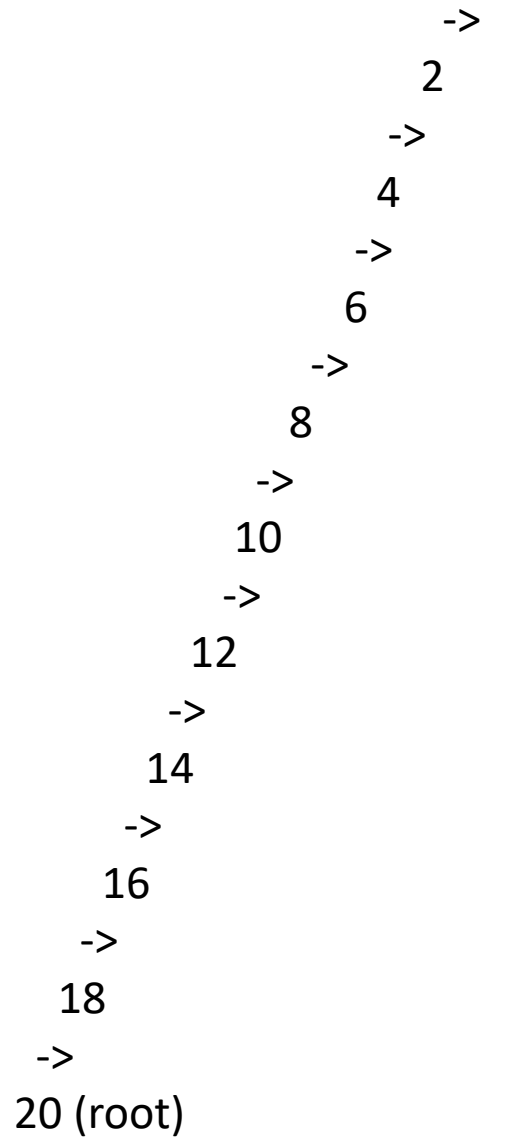
[0, 2, 4, 6, 8, 10, 12, 14, 16, 18, 20]

0 (root)

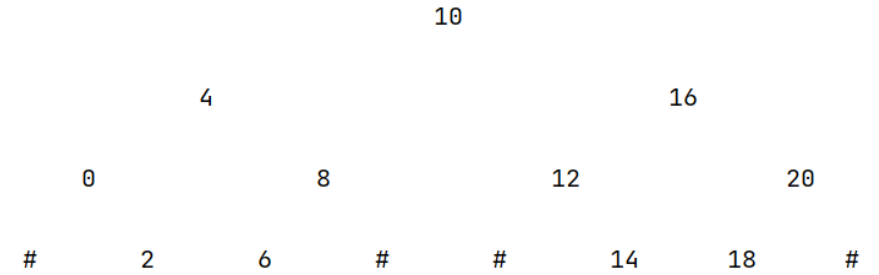


[20,18,16,14,12,10,8,6,4,2,0]

0



[10, 4, 16, 0, 2, 8, 6, 12, 14, 20, 18]



[0, 2, 4, 6, 8, 10, 12, 14, 16, 18, 20]

0 (root)

->

2

->

4

->

6

->

8

->

10

->

12

->

14

->

16

->

18

->

20

Depth 11

[20,18,16,14,12,10,8,6,4,2,0]

0

->

2

->

4

->

6

->

8

->

10

->

12

->

14

->

16

->

18

->

20 (root)

Depth 11

[10, 4, 16, 0, 2, 8, 6, 12, 14, 20, 18]

0

#

2

6

#

#

14

18

#

10

4

16

Depth 4

$\log_2 11 = 3.459 \dots$

Onward to ... optimization and profiling.

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